

COVERING THE HYPERCUBE WITH A BOUNDED NUMBER OF
DISJOINT SNAKES

JERZY WOJCIECHOWSKI

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We present a construction of an induced cycle in the n -dimensional hypercube $I[n]$ ($n \geq 2$), and a subgroup \mathfrak{H}_n of $I[n]$ considered as the group \mathbb{Z}_2^n , such that $|\mathfrak{H}_n| \leq 16$ and the induced cycle uses exactly one element of every coset of \mathfrak{H}_n . This proves that for any $n \geq 2$ the vertices of $I[n]$ can be covered using at most 16 vertex-disjoint induced cycles.

1. Introduction

Given a positive integer d , let $I[d]$ be the d -dimensional cube (also called hypercube) *i.e.* the graph with all d -tuples of binary digits as vertices, and all pairs of vertices differing at exactly one coordinate as edges. It will be convenient to think of the set of vertices of $I[d]$ as the set of elements of the group $\mathbb{Z}_2^d = (\mathbb{Z}/2\mathbb{Z})^d$ with the operation \oplus of componentwise addition mod 2. Thus, when we refer to $I[d]$, we assume that it has both a structure of a graph and a structure of a group.

A *snake* is an induced cycle in $I[d]$. For each $d \geq 2$, let $S(d)$ denote the length of the longest snake in $I[d]$. An extensive literature has evolved concerning the problem of estimating $S(d)$. See [1], [3], [5], and the references in these papers. What we know now is that

$$\frac{77}{256}2^d \leq S(d) \leq 2^{d-1} - \frac{2^{d-1}}{20d-41},$$

assuming that $d \geq 12$ in the upper bound. The lower bound was proved by Abbott and Katchalski [2] and the upper bound by Snevily [4].

During the XXIII Southeastern International Conference, Boca Raton 1992, Erdős posed the problem of deciding whether there is a number k such that for every $d \geq 2$ the vertices of $I[d]$ can be covered using at most k snakes, and if the answer to the above problem is positive, then whether it can be done in such a way that the snakes are pairwise vertex-disjoint. In this note we show that the answer to both of the above questions is positive with $k = 16$. Actually, we prove the following theorem, which is a stronger result.

Theorem 1. *For every $n \geq 2$, there is a subgroup $\mathfrak{H}_n \subset I[n]$ and a snake $C_n \subset I[n]$ such that $|\mathfrak{H}_n| \leq 16$ and C_n uses exactly one element of every coset of \mathfrak{H}_n .*

2. Basic definitions

Let $\psi_0, \psi_1: I[d] \rightarrow I[d+1]$ be the embeddings defined by

$$\psi_i(v_1 v_2 \dots v_d) = v_1 v_2 \dots v_d i,$$

for $i=0,1$. If F is a subgraph of $I[d]$, let $F^{(i)}$ be the subgraph of $I[d+1]$ obtained as the image of F under the embedding ψ_i .

For each $d \geq 2$ we define a function $H_d: \{1, 2, \dots, 2^d\} \rightarrow V(I[d])$ such that $H_d^* = (H_d(1), \dots, H_d(2^d), H_d(1))$ is a Hamiltonian cycle in $I[d]$. We set

$$H_2^* = (00, 01, 11, 10, 00),$$

and

$$H_{d+1}(i) = \begin{cases} (H_d(i))^{(0)} & \text{if } 1 \leq i \leq 2^d, \\ (H_d \circ R_d(i))^{(1)} & \text{if } 2^d + 1 \leq i \leq 2^{d+1}, \end{cases}$$

where $R_d: \{2^d + 1, 2^d + 2, \dots, 2^{d+1}\} \rightarrow \{1, 2, \dots, 2^d\}$ is the order reversing bijection,

$$R_d(i) = 2^{d+1} + 1 - i.$$

In other words, H_{d+1}^* is obtained by taking $H_d^{*(0)}$ and $H_d^{*(1)}$, removing the edges connecting their last vertices with their first vertices, joining the first vertex of $H_d^{*(0)}$ with the first vertex of $H_d^{*(1)}$ and analogously the last with the last.

Let us regard $I[d+6]$ as $I[d] \times I[6]$, that is as the d -dimensional cube $I[d]$ with each vertex being a copy of $I[6]$. Suppose that $P_d = \{(v_j^1, \dots, v_j^{r_j})\}_{j=1}^{2^d}$ is a sequence of 2^d paths in $I[6]$ such that $v_i^{r_i} = v_j^1$ when $1 \leq i \leq 2^d - 1$ and $j = i + 1$ or $i = 2^d$ and $j = 1$. Such a sequence will be called a 2^d -chain of paths in $I[6]$. We can use P_d and H_d^* to construct a cycle C_{d+6} in $I[d+6]$. Let us take the j th path $(v_j^1, \dots, v_j^{r_j})$ in the copy of $I[6]$ corresponding to the vertex $H_d(j)$ in $I[d]$; see Figure 1 for the case $d=3$.

Then, let us join $v_i^{r_i}$ from the i th copy of $I[6]$ with v_{i+1}^1 from the $(i+1)$ st copy of $I[6]$ for all $i \in \{1, \dots, 2^d\}$, where the indices are understood circularly. Hence we have

$$C_{d+6} = ((H_d(1), v_1^1), (H_d(1), v_1^2), \dots, (H_d(1), v_1^{r_1}), (H_d(2), v_2^1), \dots, (H_d(2), v_2^{r_2}), \dots, (H_d(2^d), v_{2^d}^1), \dots, (H_d(2^d), v_{2^d}^{r_{2^d}}), (H_d(1), v_1^1)).$$

It is clear that C_{d+6} is a cycle in $I[d+6]$, we will call it the cycle *generated* by H_d and P_d . We call P_d *well separated* with respect to H_d if the cycle C_{d+6} generated by H_d and P_d is an induced cycle.

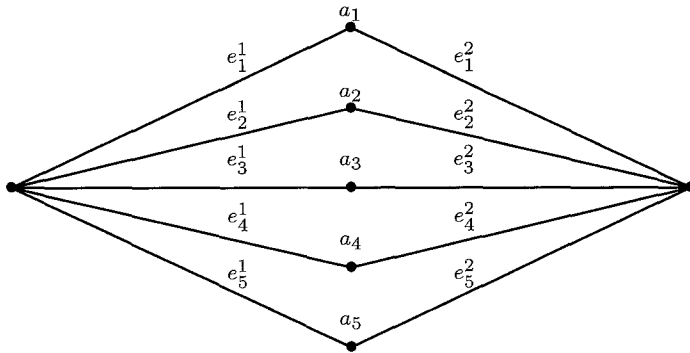


Fig. 1.

Let \mathfrak{P} be the set of all 12 paths in $I[6]$ that can be obtained from one of the following two paths

$$p_0 = (000000, 100000, 110000, 111000)$$

$$q_0 = (111000, 111100, 111110, 111111)$$

by a cyclic permutation of coordinates of every vertex of the path. Let \mathfrak{H} be the subgroup of $I[6]$ containing all $a_1 a_2 \dots a_6 \in I[6]$ such that the following two conditions are satisfied

- (i) $a_1 + a_3 + a_5$ is even, and
- (ii) $a_2 + a_4 + a_6$ is even.

Note that \mathfrak{H} is generated by the set of elements of $I[6]$ that can be obtained from 101000 by cyclic permutations of the coordinates. Also note that the elements of \mathfrak{H} are 000000, six cyclic permutations of 101000, six cyclic permutations of 111100, and three cyclic permutations of 110110, so $|\mathfrak{H}| = 16$. To prove Theorem 1, we will define a 2^d -chain P_d of paths which is well separated with respect to H_d and the paths in P_d are elements of \mathfrak{P} . We will need the following lemma.

Lemma 1. Every path $P \in \mathfrak{P}$ uses exactly one element of each coset of \mathfrak{H} in $I[6]$.

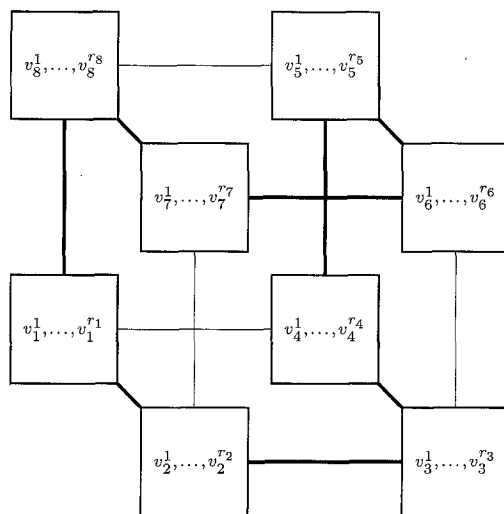
Proof. Let $P \in \mathfrak{P}$, and let v_1, v_2 be two vertices of P . It easy to observe that if $v_1 \neq v_2$ then $v_1 \oplus v_2$ has an odd number of ones or two consecutive ones (in the cyclic order), where \oplus is the operation of the group $I[6]$. Thus, if $v_1 \oplus v_2 \in \mathfrak{H}$ then $v_1 = v_2$.

This proves that P uses at most one element of every coset of \mathfrak{H} . Since $|P| \times |\mathfrak{H}| = |I[6]|$, the path P uses exactly one element of every coset of \mathfrak{H} , and the proof is complete. \blacksquare

3. The main result

Let $G = K_{2,5}$ be the graph shown in Figure 2.

The following lemma is proved in Wojciechowski [5] (Lemma 3).

Fig. 2. The graph G

- Lemma 2.** For every $d \geq 2$ there is a function $\Phi_d: \{1, 2, \dots, 2^d\} \rightarrow E(G)$ such that
- (i) if $1 \leq i \leq 2^d - 1$ and $j = i + 1$, or else $i = 2^d$, $j = 1$, then $\Phi_d(i)$ and $\Phi_d(j)$ have exactly one vertex v_i in common, such that $v_i \in \{a_1, \dots, a_5\}$ for i even, and $v_i \in \{b_1, b_2\}$ for i odd, and
 - (ii) if $(H_d(i), H_d(j)) \in E(I[d]) \setminus E(H_d^*)$, then $\Phi_d(i)$ and $\Phi_d(j)$ are vertex-disjoint. ■

Let \mathfrak{P} be the set of paths defined in Section 2. The following lemma will be used in the proof of the main result.

Lemma 3. For every $d \geq 2$ there is a 2^d -chain P_d of paths in $I[6]$ such that every path in P_d belongs to \mathfrak{P} , and P_d is well separated with respect to H_d .

Proof. Let G' be the subdivision of G obtained by subdividing e_k^1 ($1 \leq k \leq 5$) with two new vertices c_k^1 and c_k^2 in such a way that we get the path (b_1, c_k^1, c_k^2, a_k) , and subdividing e_k^2 with two new vertices c_k^4, c_k^5 , giving rise to the path (a_k, c_k^4, c_k^5, b_2) . Let $c_k^3 = a_k$ and $\xi: V[G'] \rightarrow V(I[6])$ be defined as follows. Set

$$\begin{aligned}
 \xi(b_1) &= 000000, \\
 \xi(c_1^1) &= 100000, \\
 \xi(c_1^2) &= 110000, \\
 \xi(c_1^3) &= 111000, \\
 \xi(c_1^4) &= 111100, \\
 \xi(c_1^5) &= 111110, \\
 \xi(b_2) &= 111111,
 \end{aligned}$$

and if $\xi(c_1^i) = \alpha_1 \dots \alpha_6$, then let

$$\xi(c_k^i) = \alpha_k \dots \alpha_6 \alpha_1 \dots \alpha_{k-1},$$

for $2 \leq k \leq 5$.

It is clear that the function ξ defines an embedding of G' into $I[6]$ such that the image of the subdivision of any edge of G is an induced path in $I[6]$. Let $\Phi_d: \{1, 2, \dots, 2^d\} \rightarrow E(G)$ be a function satisfying conditions (i) and (ii) of Lemma 2. Let $P_d = \{(v_i^1, \dots, v_i^4)\}_{i=1}^{2^d}$ be a 2^d -chain of paths in $I[6]$ such that (v_i^1, \dots, v_i^4) is the image under ξ of the subdivision of the edge $\Phi_d(i)$.

Let C_{d+6} be the cycle generated by H_d and P_d . Since each pair of vertex-disjoint edges in G corresponds to a pair of vertex-disjoint paths in P_d , and since each pair of edges having exactly one vertex in common corresponds to a pair of paths in P_d having exactly one vertex in common, it follows from (i) and (ii) of Lemma 2 that C_{d+6} is an induced cycle. So P_d is well separated with respect to H_d , and the proof is complete. ■

We can now prove our main result.

Proof of Theorem 1. Let us assume first that $n \geq 8$. Set $d = n - 6$. Let P_d be a 2^d -chain of paths in $I[6]$ such that every path in P_d belongs to \mathfrak{P} and P_d is well separated with respect to H_d , and let C_n be the snake generated by P_d and H_d . Let \mathfrak{H}_n be the subgroup of $I[n]$ defined to be the set of all $(a_1, a_2, \dots, a_d, b_1, \dots, b_6) \in I[n]$ such that the following conditions are satisfied.

- (i) $a_1 = a_2 = \dots = a_d = 0$,
- (ii) $b_1 + b_3 + b_5$ is even, and
- (iii) $b_2 + b_4 + b_6$ is even.

We will show that C_n contains exactly one element of every coset of \mathfrak{H}_n . Since $|C_n| = 4 \times 2^d$ and $|\mathfrak{H}_n| = 16$, it is enough to prove that C_n contains at most one element of every coset of \mathfrak{H}_n . Suppose $w_1, w_2 \in C_n$ and the cosets of w_1 and w_2 are equal. Then $w_1 \oplus w_2 \in \mathfrak{H}_n$. Let $v_1 \in I[6]$ be the sequence of the last six digits of w_1 , and similarly let $v_2 \in I[6]$ consist of the last six digits of w_2 . Since $w_1 \oplus w_2 \in \mathfrak{H}_n$, the first d digits of w_1 are the same as the first d digits of w_2 , and $v_1 \oplus v_2 \in \mathfrak{H}$. Thus v_1 and v_2 are vertices of the same path in \mathfrak{P} , and so it follows from Lemma 1 that $v_1 = v_2$. Hence $w_1 = w_2$, and the proof of the case $n \geq 8$ is complete.

If $2 \leq n \leq 6$, then we can take $C_n = (000 \dots 0, 100 \dots 0, 110 \dots 0, 010 \dots 0, 000 \dots 0)$, and \mathfrak{H}_n to be the set of all elements of $I[n]$ with first two coordinates equal 0. It is clear that C_n uses exactly one element of every coset of \mathfrak{H}_n . In the remaining case $n = 7$, let $C_7 = (0000000, 1000000, 1100000, 1110000, 1111000, 0111000, 0011000, 0001000, 0000000)$, and let \mathfrak{H}_7 be the set of all elements of $I[7]$ with the first four coordinates being either 0000 or 1010. It is straightforward to check that in this case the conclusion is true as well. Thus the proof is complete. ■

The following corollary gives the answer to the problem of Erdős.

Corollary 1. For every $n \geq 2$ the vertices of $I[n]$ can be covered using at most 16 pairwise disjoint snakes.

Proof. Let us take a subgroup \mathfrak{H}_n and a snake C_n in $I[n]$ such that $|\mathfrak{H}_n| \leq 16$ and C_n uses exactly one element of every coset of \mathfrak{H}_n . The family of snakes $\{C_n \oplus h: h \in \mathfrak{H}_n\}$ contains 16 vertex-disjoint snakes covering $I[n]$. ■

4. Concluding remarks

Let k_0 be the smallest integer such that for every $n \geq 2$, the cube $I[n]$ can be vertex covered by at most k_0 snakes. Let k_1 and k_2 be defined in a similar way taking pairwise vertex-disjoint snakes and pairwise vertex-disjoint snakes of equal length, respectively. Set k_3 to be the smallest integer such that for every $n \geq 2$ there is a subgroup \mathfrak{H}_n and a snake C_n in $I[n]$ such that $|\mathfrak{H}_n| \leq k_3$ and C_n uses exactly one element of every coset of \mathfrak{H}_n . As a corollary of Theorem 1 and the upper bound for the length of snakes we get the following theorem.

Theorem 2. We have $3 \leq k_0 \leq k_1 \leq k_2 \leq k_3 \leq 16$ and $k_2, k_3 \in \{4, 8, 16\}$. ■

The question of determining the exact values of k_0 , k_1 , k_2 , and k_3 remains open. It might be possible to modify the technique used in this note to improve Theorem 2 to make the upper bound equal 8 or perhaps even 4. The possible approach may involve finding a more sophisticated embedding of a subdivision of $K_{2,5}$ into $I[c]$, where c is a small integer constant (in our proof we used $c=6$), or else replacing $K_{2,5}$ by another graph, perhaps $K_{2,4}$.

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Jerzy Wojciechowski

Department of Mathematics
West Virginia University
PO BOX 6310, Morgantown
WV 26506-6310, USA
 UNO2O243@WVNVMS.WVNET.EDU
 JERZY@MATH.WVU.WVNET.EDU